

Privacy-aware traffic monitoring In Cloud Computing

Nikhil adgokar Amol v zade

Abstract— Cloud security becomes one of the major barriers of a widespread adoption of conventional cloud services. anticipate that the same problems will be present in VCs. In a VC, underutilized vehicular resources including computing power, storage, and Internet connectivity can be shared between drivers or rented out over the Internet to various customers. Clearly, if the VC concept is to see a wide adoption and to have significant societal impact, security and privacy issues need to be addressed. The main contribution of this work is to identify and analyze a number of security challenges and potential privacy threats in VCs. Although security issues have received attention in cloud computing and vehicular networks, we identify security challenges that are specific to VCs, e.g., challenges of authentication of high-mobility vehicles, scalability and single interface, tangled identities and locations, and the complexity of establishing trust relationships among multiple players caused by intermittent short-range communications. Additionally, we provide a security scheme that addresses several of the challenges discussed. Vehicles often communicate through multihop routing. A request response will include multiple participants, including users, infrastructure, servers, platform, application, and key generator and privacy agent.

Index Terms:- Challenge analysis, cloud computing, privacy, security, vehicular cloud.

1. INTRODUCTION

In an effort to help their vehicles compete in the marketplace, car and truck manufacturers are offering increasingly more potent onboard devices, including powerful computers, a large array of sensors, radar devices, cameras, and wireless transceivers. These devices cater to a set of customers that expect their vehicles to provide seamless extension of their home environment populated by sophisticated entertainment centers, access to Internet, and other similar wants and needs. Powerful onboard devices support new applications, including location-specific services, online gaming, and various forms of mobile infotainment. In which computer processing is performed in the internet "cloud," this means that user need not concern themselves with the processing details. Although cloud computing enables flexible and agile computing impossible with existing system, it brings new security problems that may users anxious about safety and reliability. This paper describes the security problem surrounding cloud computing and present existing approaches to solving them. It also described the security architectures of a service platform proposed by Fujitsu for dealing with those problems. Superficially, the security issues encountered in VCs may look deceptively similar to those experienced in other networks. However, a more careful analysis reveals that many of the classic security challenges are exacerbated by the characteristic features of VCs to the point where they can be construed as VC-specific.

For example, the high mobility of vehicles is apt to cause significant challenges related to managing authentication, authorization, and accountability since the vehicles communicate through short-range dedicated short-range communications (DSRC) transceivers. Vehicular mobility and tangled identities and locations also cause significant challenges of privacy. Employing pseudonyms is a common solution, but the high mobility makes the task of updating pseudonyms quite difficult. The two main contributions of this work are to identify and analyze security challenges and privacy threats that are VC specific and to propose a reasonable security framework that addresses some of the VC challenges identified in this paper.

2. RELATED WORKS

Although handling security issues in VANET is very tough, because handling security issues will increase the overhead cost and also the functional cost. VANET will be executed when cost management and security handling issues, both will be reduced or compromised so that the system becomes effective from both the point of views. While going through all the papers each and every paper gave us certain information [1]. VANET follows a simple security architecture which is underlined below. The basic architecture consists of Network nodes which can be either Vehicles or Road Side Infrastructure and existing Registration Authorities for vehicle registration and record maintenance. These nodes will be installed with required sensors for gaining information, processing units for processing the collected or received information and communication system for disseminating information to and receiving

• R.B.Garad is currently pursuing bachelors degree in Computer Science & engineering in SGBA University, Amravati, India, 444709. E-mail: adgokar.nikhil@gmail.com

ing information from other nodes [1].

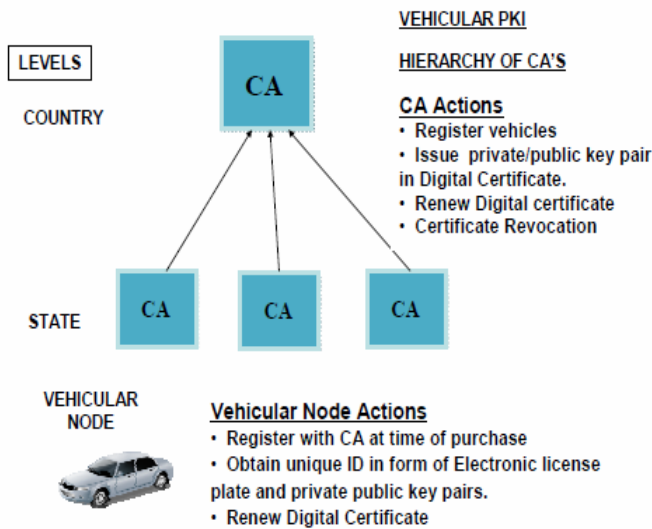


Fig 2.1 VANET security architecture

The basic architecture consists of Network nodes which can be either Vehicles or Road Side Infrastructure and existing Registration Authorities for vehicle registration and record maintenance. These nodes will be installed with required sensors for gaining information, processing units for processing the collected or received information and communication system for disseminating information to and receiving information from other nodes [1]. A secure system, besides the basic network nodes, will consist of a Vehicular Public Key infrastructure (PKI), a Secure Computing platform and various security mechanisms. Secure mechanisms comprise identity management using Electronic License Plates with certified public and private keys attached to the owner, Authentication and Integrity using Digital Signatures, Privacy using Pseudonyms, Pseudonym handling and Certification Revocation mechanisms.

A Vehicular PKI will consist of the national and state level registration authorities acting as Certification Authorities (CAs) which will issue certified public/private key pairs to vehicles. A Secure Computing platform on a vehicle will consist of tamper resistant hardware and firmware. Its job is to store cryptographic material (private keys) and a trusted (tamper proof) clock [1].

Digital Signatures will provide the required authentication and integrity along with non-repudiation using timestamps. Privacy is introduced by using Pseudonyms in the form of additional set of public/private keys which are given to the user. These keys are used for a short period of time and changed frequently. These keys do not contain identity related information but can be traced back to the owner in liability related cases with the help of central authorities. The aim in using pseudonyms is to ensure that a vehicle cannot be tracked and a message cannot be attributed to its sender by other vehicles. Finally, when a vehicle becomes faulty or is

detected as an illegitimate or malicious vehicle, Certificate Revocation mechanisms are required to revoke both long term certificates and set of Pseudonyms currently being used by the vehicle.

The security architecture developed by the Vehicle Safety Communications Consortium (VSCC) and subsequently submitted. The only approach for a security architecture in vehicular networks that is under standardization so far. It defines a public-key-infrastructure (PKI)-based approach for securing messages sent in a vehicle-to-vehicle and vehicle-to-infrastructure fashion. The Daimler Chrysler group also published security architecture in the form of a layered structure with multiple views of the system. The security architecture of the system discussed in this paper contains the Vehicle Manufacturer and the Registration Authority for registration of nodes and assigning node identifiers, the Inspection site for test and certification of nodes, an Escrow entity with authority to identify and revoke certification of nodes and finally the communication infrastructure consisting of communication systems, processing and databases necessary to carry out online testing, pseudonym provision for nodes and infrastructure based data assessment and intrusion handling. There are also some papers which dealt with the entire environment that how the communication process will be like, its result showed that effective communication between nodes depends on the density of vehicular nodes, their velocities and the number of lanes, i.e. width of the road. Papers such as, dealt with routing protocols and gave effective solutions so that the communication between the nodes is computationally effective and leading to less congestion of network traffic.

There are security solutions that are related with the deployment of RSU's. In CA cluster in different regions comply with corresponding scalability strategy and regional policy. A distributed IDS system integrated with the CA database provide further security protection from malicious vehicles with legal certificates. The certificate caching and forwarding schema accelerates authentication. Where as in usage of DSRC mainly gives a flawed solution in deployment of RSU, but it gave a simple mathematical approach of getting the position of a vehicular node without the help of GPS.

In a solution of group formation combined with RSU is illustrated, which resulted in easy revocation of malicious vehicle, location privacy protection is improved and the system maintenance becomes flexible. In, the seminar has made use of syntactic aggregation and cryptographic aggregation techniques to dramatically reduce the transmission cost, and adopt batch verification technique for efficient emergency messages verification. This made the authentication of the emergency events easier.

3. SECURITY CHALLENGES IN

VEHICULAR COMMUNICATION

The security challenges in VC are a new, exciting, and unexplored topic. Vehicles will be autonomously pooled to create a cloud that can provide services to authorized users. This cloud can provide real-time services, such as mobile analytic laboratories, intelligent transportation systems, smart cities, and smart electric power grids. Vehicles will share the capability of computing power, Internet access, and storage to form conventional clouds. These researchers have only focused on providing a framework for VC computing, but as already mentioned, the issue of security and privacy has not yet been addressed in the literature. cloud security becomes one of the major barriers of a widespread adoption of conventional cloud services[2]. we anticipate that the same problems will be present in VCs. Recently, vehicular ad hoc network (VANET) security and privacy have been addressed by a large number of papers. Algorithms Radar can be employed as a "virtual eye," and onboard radar can detect the location of vehicles. Public Key Infrastructure (PKI) and digital signature-based methods have been well explored in VANETs [3]. A certificate authority (CA) generates public and private keys for nodes. The purpose of digital signature is to validate and authenticate the sender. The purpose of encryption is to disclose the content of messages only to entitled users. PKI is a method that is well suited for security purposes, particularly for roadside infrastructure [3]. Geo Encrypt in VANETs has been proposed by Yan *et al.*. Their idea is to use the geographic location of a vehicle to generate a secret key. Messages are encrypted with the secret key, and the encoded texts are sent to receiving vehicles [4]. The receiving vehicles must be physically present in a certain geographic region specified by the sender to be able to decrypt the message. Recently, some attention has been devoted to the general security problem in clouds, although not associated with vehicular networks. The simple solution is to restrict access to the cloud hardware facilities. A trust coordinator maintained by an external third party is imported to validate the entrusted cloud manager, which makes a set of virtual machines (VMs) such as Amazon's E2C (i.e., Infrastructure as a Service, IaaS) available to users. Garfinkel proposed a solution to prevent the owner of a physical host from accessing and interfering with the services on the host adopted a similar solution. When a VM boots up, system information such as the basic input output system (BIOS), system programs, and all the service applications is recorded, and a hash value is generated and transmitted to a third-party Trust Center. For every period of time, the system will collect system information of the BIOS, system programs, and all the service applications and transmit the hash value of system information to the third-party Trust Center. The Trust Center can evaluate the trust value of the cloud. Krautheim also proposed a third party to share the responsibility of security in cloud computing between the service provider and client, de-

creasing the risk exposure to both. Jensen [4]. stated technical security issues of using cloud services on the Internet access. Wang [4]. proposed public-key-based homomorphic authenticator and random masking to secure cloud data and preserve privacy of public cloud data. The bilinear aggregate signature has been extended to simultaneously audit multiple users. Ristenpart presented experiments of locating co-residence of other users in cloud VMs.

4. VEHICULAR CLOUDS

4.1. Conceptual Overview

4.1.1 Cloud Computing

In recent years, cloud computing and its myriad applications that promise to change the way I think about computing and data storage have received a huge amount of attention. Cloud users do not need to install expensive hardware and software on their local machine. They can subscribe and use both hardware and software *as a service* when they want to use it. In addition, fees are charged based on the usage of the service. The users can access these services through Internet browsers, and no expensive client terminals are needed. Service providers can make good use of *excess* capabilities on the server side including processors, storage, and sensors that can be used to provide services to clients [23].

4.1.2 VANET

In VANETs, the vehicles communicate with each other and/or with the roadside infrastructure using the Federal Communications Commission-mandated DSRC [24], restricting the transmission range to 300–1000 m. There are two types of VANET networks: the zero-infrastructure and the infrastructure-based VANET. The zero-infrastructure VANET is created on-the-fly. There are many challenging security and privacy problems because no infrastructure is used for authentication and authorization. The infrastructure-based VANET can be formed based on the roadside infrastructure. The infrastructure can act as wireless access points for authentication and authorization purposes [24]. By the same token, the vehicles can use the infrastructure to report events and to exchange information.

4.1.3 Vehicular Clouds (VCs)

Similar to VANETs, there are two types of VCs. In the first type called Infrastructure-based VC, drivers will be able to access services by network communications involving the roadside infrastructure. In the second type called Autonomous VC (AVC), vehicles can be organized on-the-fly to form VC in support of emergencies and other ad hoc events [2]. VCs provide services at three levels, i.e., application, platform, and infrastructure. Service providers use the levels differently based on what and how the services are offered. The fundamental level is called Infrastructure as a Service (IaaS) where infrastructure such as computing, storage, sensing, communicating devices, and software are created as VMs. The next level is Platform as a Service (PaaS), where components and ser-

vices (such as httpd, ftpd, and email server) are provided and configured as a service. The top level is called Software as a Service (SaaS), where applications are provided in a “pay-as-you-go” fashion. VCs provide a cost-efficient way to offer comprehensive services. For example, a cheaper vehicle with network access can access a VM with strong computation, communication, sensing capability, and large storage. Many applications such as traffic news, road conditions, or intelligent navigation systems can be provided by a VM [25].

4.2. Potential Applications of VC Computing

In this section, we review several possible applications of VCs.

- **Vehicle maintenance:** Vehicles receive software updates from cloud whenever vehicle manufacturers upload a new version of software.
- **Traffic management:** Drivers can receive traffic status reports (e.g., congestion) from VCs.
- **Road condition sharing:** Road conditions such as flooding areas and black ice on the roadway can be shared in VCs. Drivers will be alerted if there are serious road conditions.
- **Accident alerts at intersections:** Under demanding driving conditions such as fog, heavy storm, snow, and black ice, drivers can order this service to alert them of possible accidents at intersections. Infrastructure, e.g., a tall building, can include high-precision radar to detect car accidents. This infrastructure will cover the whole intersection and frequently scan the intersection. An intelligent algorithm will be applied to each scan result to predict the possibility of accidents.
- **Safety applications:** Applications related to life-critical scenarios such as collision avoidance and adaptive cruise control require strong security protection, even from surrounding environmental security threats.
- **Intelligent parking management:** Vehicles will be able to book a parking spot using the VC. All the parking information will be available on clouds without central control. Requests from different physical places can be transferred to the most desired parking lots.
- **Planned evacuations:** In some disasters such as a hurricanes and tsunamis, VCs will be instrumental in organized evacuations.

5. ANALYZING SECURITY IN A VEHICULAR CLOUD

In this section, we introduce a set of security analyses that are specially associated with VCs.

5.1. Security and Privacy Attacks in VC

5.1.1. Attacker Model

Traditional security systems are often designed to prevent

attackers from entering the system. However, security systems in the VC have a much harder time keeping attackers at bay, because multiple service users with high mobility can share the same physical infrastructure. In the VC environment, an attacker can equally share the same physical machine/infrastructure as their targets, although both of them are assigned to different VMs. To this point, attackers can have more advantages than the attackers on traditional systems. In addition, the attackers are physically moving from place to place as vehicles are mobile nodes. It is much harder to locate the attackers. The main targets of an attacker are given as follows:

- **Confidentiality**, such as identities of other users, valuable data and documents stored on the VC, and the location of the VMs, where the target's services are executing;
- **Integrity**, such as valuable data and documents stored on the VC, executable code, and result on the VC;
- **Availability**, such as physical machines and resources, privileges, services, and applications.

One possible form of attack is given below:

- 1) Find the geographic location of the target vehicle and physically move close the target machine.
- 2) Narrow down the possible areas where the target user's services are executing by mapping the topology of VC.
- 3) Launch multiple experimental accesses to the cloud, and find out if the target user is currently on the same VM.
- 4) Request the services on the same VM where the target user is on.
- 5) Use system leakage to obtain higher privilege to collect the assets [23].

Due to the features of the VC, there are several challenges for attackers as well. High mobility of vehicles is like a double-edged sword. It makes it hard for attackers to harm a specific target vehicle. First, the vehicle's access of each virtual machine can be transitory as vehicles constantly move from one district to another one, if each district is associated with a virtual machine. Additionally, attackers need to locate on which machine/infrastructure a specific target is located because all users in the VC are distributed on virtual machines. However, it is possible to locate the co-residence of other users. Experiments have been done to catch and compare the memory of processors, and users can find co-residence in the same physical machine. Third, the attackers must be physically co-located with the target user on the same physical machines [23]. This will require attackers to be physically present at the same region with the target vehicles or shadow with the target vehicles at the same speed. These challenges make attacking extremely difficult because coexistence is hard to achieve and is

temporary. Finally, the attackers have to collect valuable information with certain privileges or with security tokens.

5.1.2. Threats

The threats in the VC can be classified using STRIDE, a system developed by Microsoft for classifying computer security threats. The threat categories are given here.

- **Spoofing user identity:** The attackers pretend to be another user to obtain data and illegitimate advantages. One classic example is the “man-in-the-middle attack,” in which the attackers pretend to be Bob when communicating with Alice and pretend to be Alice when communicating with Bob. Both Alice and Bob will send decrypt table messages to the attackers.
- **Tampering:** The attackers alter data and modify and forge information.
- **Repudiation:** The attackers manipulate or forge the identification of new data, actions, and operations.
- **Information disclosure:** The attackers uncover personally identifiable information such as identities, medical, legality, finance, political, residence and geographic records, biological traits, and ethnicity.
- **Denial of Service:** The attackers mount attacks that consume system resources and make the resources unavailable to the intended users.
- **Elevation of privilege:** The attackers exploit a bug, system leakage, design flaw, or configuration mistake in an operating system or software application to obtain elevated access privilege to protected resources or data that are normally protected from normal users.

5.2 Authentication of High-Mobility Nodes

Security authentication in the VC includes verifying user identity and message integrity. To conduct authentication, there are some metrics that can be adopted [27].

- **Ownership:** A user owns some unique identity (e.g., identity card, security token, and software token).
- **Knowledge:** A user knows some unique things [e.g., passwords, personal identification number and human challenge response (i.e., security questions)].
- **Biometrics:** These include the signature, face, voice, and fingerprint.

However, it is challenging to authenticate vehicles due to high mobility. First, high mobility makes it hard to authenticate messages with a location context. For example, accident alert message associated with locations and events at a specified time are hard to verify because the locations of vehicles are constantly changing. Second, high mobility and a short transmission range may result in the recipient being out of reach. It is likely that a vehicle at the border of access point can change its access point when the authentication message is transmit-

ted back. Third, the security token (security key pairs) is hard update. Some vehicles can even park for years without starting a single time. These situations will make the updating tasks of the security token significantly difficult. In addition, it is challenging to authenticate a vehicle's or driver's identity in the VC. To protect privacy, these identities are often replaced by pseudonyms. The authentication of identity can be complex and makes Sybil attacks possible [27].

5.3 Establishing Trust Relationships

Trust is one of the key factors in any secure system. A trust relationship can exist in several ways. The network service providers and the vehicle drivers have access to trust. There will be a large number of government agents, e.g., the Department of Motor Vehicles (DMV) and the Bureau of Motor Vehicles (BMV) are trusted organizations. The relationship between the BMV and vehicle drivers is identity uniqueness and legitimacy. However, the large population of vehicles creates challenges to building trust relationships to all the vehicles at any time. There will be occasional exceptions. In addition, drivers are increasingly concerned about their privacy. Tracking vehicles/drivers will cause worries in most cases. As a result, pseudonyms are often applied to vehicles. On the other hand, a certain level of trust of identity is needed. Some applications such as accident reliability investigation by law enforcement or insurance companies require the driver's identity to be responsible for accidents. Therefore, we assume that a low level of trust relationship exists in VANETs. To obtain a high-level trust relationship, the security scheme discussed in Section IV needs to be executed. In VCs, it is far more challenging to build trust relationships than in vehicular networks and conventional cloud computing. Fig. 2 shows an example of multiple participants in a VC. The VC is often based on DSRC. Many applications need multihop routing, with multiple nodes involved in communication. Therefore, the VC has inherited the challenge of establishing trust relationships among multiple vehicles, roadside infrastructure, service providers, network channels, and even the secret key generator. In this seminar, I assume that the VC cloud infrastructure is trusted, the VC service providers are trusted, the vast majority of VC users are trustworthy, and the attackers have the same privileges as normal users.

5.4 Location Validation and Pseudonymization

Most, if not all, VC applications rely on accurate location information. Therefore, location information must be validated. There are two approaches to validate location information: active and passive. Vehicles or infrastructure with radar (or camera, etc.) can perform active location validation. Radar input can be used to validate location information. Vehicles or infrastructure without radar, or in a situation where radar detection is not possible, can validate location information by applying statistical methods. A vehicle's identity is often tan-

gled with owner's identity. Because of legal and insurance issues, a vehicle's unique identity (such as vehicle identity number, Internet Protocol address, and hostname) is often linked to the owner's identity. Therefore, tracking a vehicle can often invade its owner's privacy. To protect privacy, one can replace vehicular identity by a pseudonym. The real identity can only be discovered by the Pseudonymization Service Center, which is a secured and trusted entity. The pseudonym is subject to timeout. After expiration, a new pseudonym will be assigned [9].

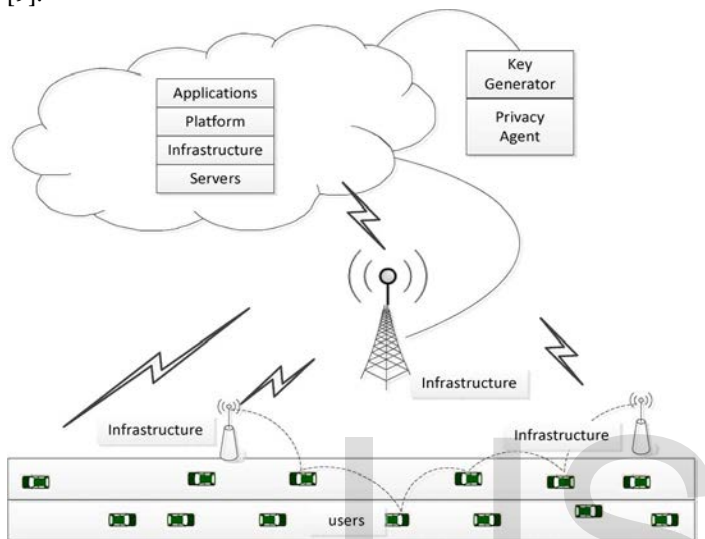


Fig.5.1 Vehicles often communicate through multihop routing. A request response will include multiple participants, including users, infrastructure, servers, platform, application, and key generator and privacy agent.

5.4 Location Validation and Pseudonymization

Most, if not all, VC applications rely on accurate location information. Therefore, location information must be validated. There are two approaches to validate location information: active and passive. Vehicles or infrastructure with radar (or camera, etc.) can perform active location validation. Radar input can be used to validate location information. Vehicles or infrastructure without radar, or in a situation where radar detection is not possible, can validate location information by applying statistical methods. A vehicle's identity is often tangled with owner's identity. Because of legal and insurance issues, a vehicle's unique identity (such as vehicle identity number, Internet Protocol address, and hostname) is often linked to the owner's identity. Therefore, tracking a vehicle can often invade its owner's privacy. To protect privacy, one can replace vehicular identity by a pseudonym. The real identity can only be discovered by the Pseudonymization Service Center, which is a secured and trusted entity. The pseudonym is subject to timeout. After expiration, a new pseudonym will be assigned [9]. plates (DLPs) or electronic license plates, which are a wireless device periodically broadcasting a unique identity string,

have been proposed. Temporary public keys as DLPs can protect privacy and can be broadcast.

5.5 Scalability

Security schemes for VCs must be scalable to handle a dynamically changing number of vehicles. Security schemes must handle not only regular traffic but special traffic as well, e.g., the large volume of traffic caused by special events (e.g., football games, air shows, etc.) The dynamics of traffic produces dynamic demands on security. For example, imagine a downtown area with several supermarkets and stores that take orders from vehicles in traffic, complete with credit card information. To protect credit card information, comprehensive cryptographic algorithms must be applied. However, the comprehensive algorithms decrease the efficiency of communication response time. Therefore, better algorithms and, perhaps, less comprehensive security schemes are needed to speed up the response time.

5.6 Single-User Interface

Single-user access interface is another challenge to VCs. When the number of service accesses in a cloud increases, the number of VMs that provide the service will increase to guarantee quality of service. More VMs will be created and assigned. With the increase in VMs, security concerns grow as well. When the number of service accesses decreases, the number of VMs that provide the service will decrease to improve resource utilization. Some VMs will be destroyed and recycled. These procedures are transparent to vehicles. Vehicles only see one access interface and do not need to know the changing of VMs. To achieve scalability, a simple solution is to clone and expand the service in a different cloud. However, a single interface obviously makes scalability even more difficult.

5.7 Heterogeneous Network Nodes

Conventional cloud computing and fixed networks often have homogeneous end users. As it turns out, vehicles have a large array of (sometimes) vastly different onboard devices. Some high-end vehicles have several advanced devices, including a Global Positioning System (GPS) receiver, one or more wireless transceivers, and onboard radar devices. In contrast, some economy models have only a wireless transceiver. Some other vehicles have different combinations of GPS receivers, wireless transceivers, and radar. Different vehicle models have different device capabilities such as speed of processor, volume of memory, and storage. These heterogeneous vehicles as network nodes create difficulties to adapting security strategies. For example, PKI encryption and decryption algorithms will require vehicles to meet certain hardware conditions.

5.8 VC Messages

5.8.1 Safety Messages:-

The initial motivation of VANET was the dissemination of traffic safety messages. Based on the emergency level, there are three types of safety messages.

- **Level one:** public traffic condition information. Vehicles exchange traffic information (e.g., traffic jam) that indirectly affects other vehicles' safety, e.g., a traffic jam increases the likelihood of accidents. This type of message is not sensitive to communication delay, but privacy needs to be protected.
- **Level two:** cooperative safety messages. Vehicles exchange messages in cooperative accident avoidance applications. These messages are often time critical, and privacy needs to be protected.
- **Level three:** liability messages. After accidents happen, there will be liability messages generated by law enforcement authorities. These messages contain important evidence for liability claims and are bonded by a certain time range. Privacy information is naturally protected.

A common format of safety messages is timestamp, geographic location, speed, percentage of speed change since the last message, direction, acceleration, and percentage of acceleration change since last message. The safety message will append information such as public traffic condition and accidents. The appended message can help determine liability.

5.9 Key Management

5.9.1 Key Assignment and Rekeying

In VANETs, some organizations can serve as CAs: governmental transportation authorities, vehicle manufacturers, or nonprofit organizations. Initially, a vehicle will receive a key pair from the manufacturer or some governmental authority. Key assignment is on the basis of a unique ID with a certain expiration time. Upon expiration, the key pair has to be renewed at the local DMV/BMV. The renewal/expiration period can be the same period of vehicular state inspection, e.g., mandatory annual state inspection in many U.S. states.

5.9.2 Key Verification

To verify key pairs, we assume that every vehicle trusts CAs and that CAs are tamper-proof. Key validation can be done at the CAs or sub-CAs. Let pub_i of vehicle i be the public key issued by a CA j , i.e., CA j . Vehicle i will have a certificate $Cert_i[pub_i]$ assigned by CA j when CA j assigns the public key. The process of validating public key will compute the following certificate at CA j . $Cert_i[pub_i] = pub_i \parallel sig_{pri_{CA_j}}(pub_i \parallel ID_{CA_j})$ where pri_{CA_j} is the private key of CA j , and ID_{CA_j} is the identity of CA j . The idea is to sign the special message $pub_i \parallel ID_{CA_j}$ using the private key of CA j . The digital signature algorithm has been discussed in Section .

5.9.3 Key Revocation

Key revocation is an important and effective way to prevent attacks. There are certain cases when key pairs will be exposed to attackers. It is obvious that an exposed key pair needs to be disabled. One of the advantages of PKI is that PKI can revoke a key pair. Vehicles will be aware that the exposed

key pair has been revoked and refuse to communicate with vehicles with invalid key pairs. PKI uses certificate revocation lists (CRLs) to revoke keys. CRLs include a list of the most recently revoked certificates and are instantly distributed to vehicles. In VANETs, the infrastructure can serve as CRL distributors [9].

The CAs can revoke key pairs by using onboard tamperproof devices. Suppose that CAs want to revoke the key pairs of vehicle v . CAs will send out the revoke message signed by public key of V to the tamper-proof devices. After receiving this revoking message, the tamper-proof device will validate the message and revoke the key pairs. The tamper-proof device will also send back an ACK to the CA to confirm the operation. To improve communication between V and CA, the vehicle's location is retrieved to select the closest CA. If the latest vehicle location failed to be retrieved, the last location will be used to select the closest CA. In this case, the CA will use a broadcasting message to revoke the key pairs. The broadcasting message can be sent out by using several media such as FM, Internet, and satellite.

To avoid attackers reporting other vehicles to CA to revoke the key pairs of other vehicles, revocation will be triggered by a certain number of neighboring vehicles. There is another risk that attackers can launch planned attacks. For example, several attackers can surround a well-behaved vehicle and report the well-behaved vehicle as a misbehaving vehicle. Prevention of this risk is very challenging. Due to the dynamics of traffic, it is costly to launch such an attack. One possible solution is to build behavior history records and credit the past behavior into values, just like the bank credit system. A similar solution has been discussed as Map History [9].

6. RESEARCH APPROACH

In this section, we offer a first attempt to addressing several of the challenges previously discussed. We begin by describing the two VC models, i.e., infrastructure- and ad-hoc-based models. We then demonstrate algorithms to enhance authentication of high-mobility vehicles, configure customized security schemes, and improve scalability of security schemes

6.1. The Cloud Model

The cloud in this proposal is associated with a number of grids. A city or a traffic area is partitioned into grids. The grid size is predefined, e.g., 700 m² and with two GPS coordinates. The grid of a city is shown in Fig. 6.1. Each cell is associated with a virtual machine in the cloud. The virtual machine can dynamically request resources from cloud. For example, when the grid is congested, the corresponding virtual machine will request more communicating, storage, and computing resources. The cloud will be able to borrow these resources from the idle virtual machine, which is associated with sparse traffic grid. Therefore, the traffic of the whole city can be

mapped to the cloud...

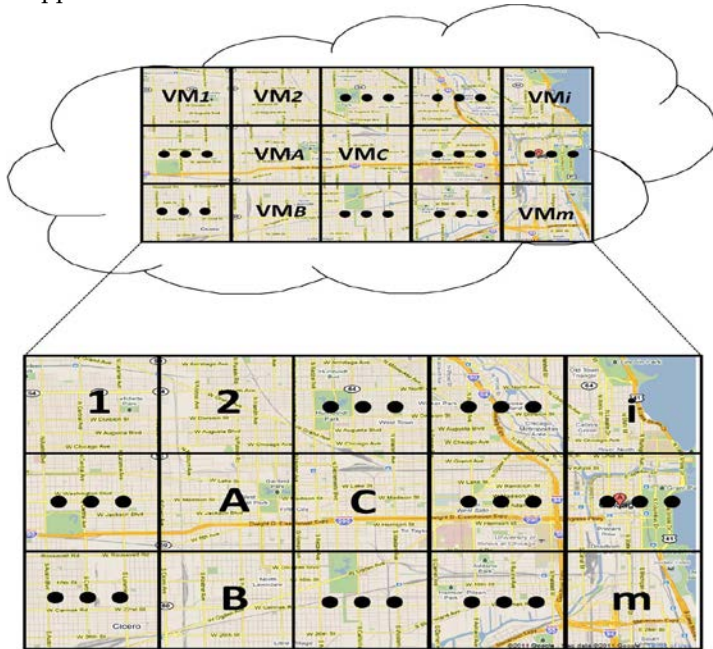


Fig.6.1 Downtown area partitioned into cells, each mapped to a virtual machine.

This cloud model provides high capability in customizing cloud services and the security scheme. For example, a downtown area is often queried about vacant parking spots and congestion status. The corresponding virtual machine can be specially configured and optimized in the smart parking and congestion control services. At a busy intersection, a collision warning service can be specialized and optimized in the virtual machine. A possible solution is to collect and sort all the vehicles' mobility information at the intersection. When vehicles are too close to each other by considering the headway distance and relative speed, the vehicles will receive an alarm from the cloud. Even cheaper cars that have no radar cruise control system can get benefits from the cloud collision warning system. What distinguishes vehicles from standard nodes in a conventional cloud is autonomy and mobility. Indeed, large numbers of vehicles spend substantial time on the road and may.

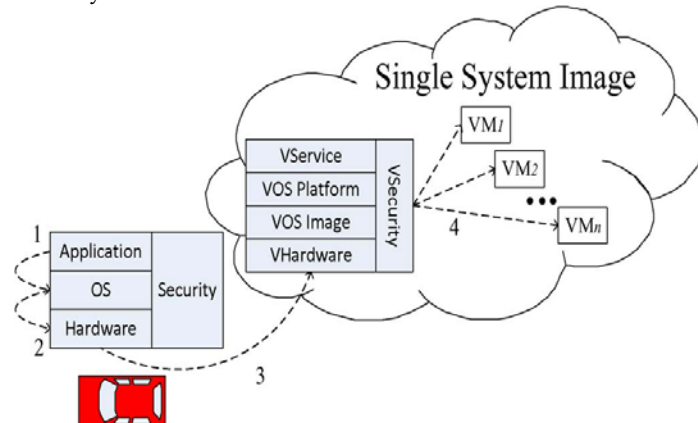


Fig.6.2 Vehicle node in a cell can communicate with a virtual

machine that is responsible for the cell.

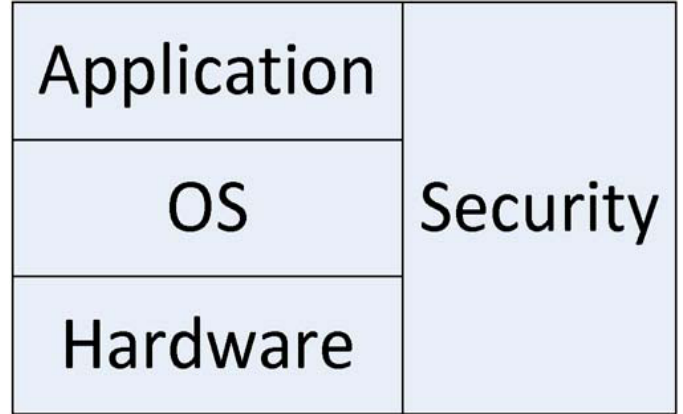


Fig.6.3 Vehicle node image is located on each individual vehicle.

Be involved in dynamically changing situations; we argue that, in such situations, the vehicles have the potential to cooperatively solve problems that would take a centralized system an inordinate amount of time, rendering the solution useless. Vehicles automatically form a cloud by connecting virtual cells, which can be a group of vehicles. Each virtual cell is associated with a virtual machine in which vehicles rent or contribute their spare computing, storage, and sensing resource. The group of vehicles moves at almost the same speed. Since vehicles are cloud constructors and cloud users, all vehicles inside a cell can directly receive packets from each other. A cell leader can be elected to communicate with other clouds [9].

6.2 Virtual Machines of VCs

This objective concerns how a cloud is formed and how the service can be provided. We first consider the basic modules of the VC and then introduce the process of a service request and response. The communication between a vehicle and the cloud is through a unique entry. The cloud provides a single system image to each individual virtual machine shown as Fig. 6.2. Each vehicle has a node image, which includes hardware drivers, operating system image, security system, and applications, as shown in Fig. 6.3. When the applications of the vehicle send a request to the cloud, the request will be forwarded to the operating system and, then, the hardware (network driver). The request will be sent by the wireless network and received by the cloud single system image. The allocator of the cloud will locate which virtual machine should be responsible for the request and forward the request to the virtual machine. If the request needs to access other virtual machines, e.g., to check the traffic congestion status of a city in a remote state, the virtual machine can communicate with other virtual machines as well.

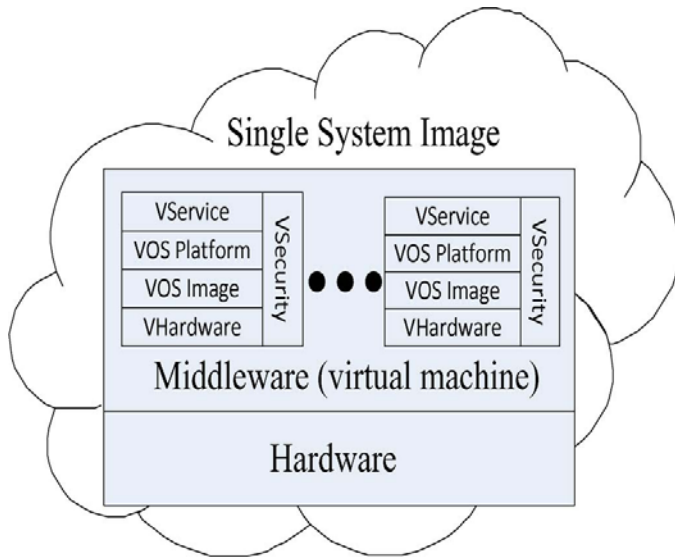


Fig.6.4 Cloud provides a single system image and is composed by a number of virtual machines.



Fig.6.5 Single virtual machine located in the cloud.

The VC is a single system image composed of a number of virtual machines. A single image can be created by a layer of middleware between the hardware manager system and a number of virtual machines, as shown in Fig. 6.4. The middleware is a cloud operating system and a platform to allocate a large number of virtual machines. Each virtual machine is composed of virtual hardware, virtual operating system image, virtual operating system platform, virtual security system, and virtual services, as shown in Fig. 6.5. The virtual hardware is composed of several real computers that virtually act as real hardware and provide the interface of the hardware. The virtual operating system image can be any current operating system, such as Linux/Unix or Windows. The virtual operating system platform includes not only the operating system but system applications such as web server and databases. The virtual security system is a set of complete security solutions, including hardware and software. The customized security protocols can be configured and replaced in this module. The virtual services are actual services that are con-

figured for the related traffic area/grid.

6.3 Securing VCs

6.3.1 Trust Relationship

For infrastructure-based VC, trust relationships can be built by infrastructures that are constructed by authorities such as BMV/DMV or other transportation agencies. Infrastructure will be authenticated and assigned with security key pairs. Infrastructure stores the key pairs in tamperproof devices. As shown in Fig. 2, vehicles communicate with.

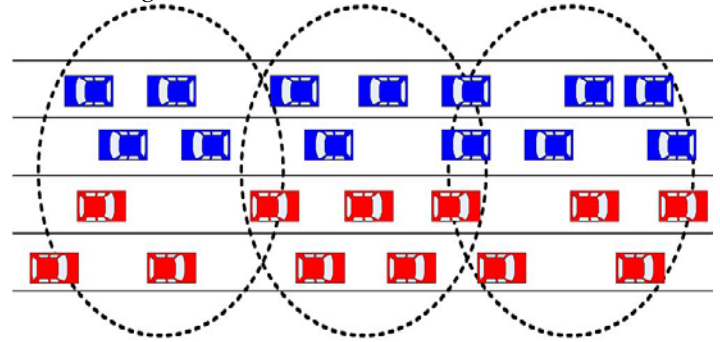


Fig.6.6 Trust relationship in AVCs can be built on the basis of a group of vehicles. The behavior of a vehicle can be monitored by all members.



Fig.6.7 Geographic location-based security mechanism.

The shaded square is the naval base. Only the vehicles in the shaded rectangle region (i.e., vehicle g can decrypt and access the received cipher text sent by vehicle a). infrastructure as access point to the VC. The infrastructure is sufficiently capable to handle large numbers of accesses in its transmission range. The scalability of trust relationships can be achieved because the infrastructure is connected to each other by fixed networks. For AVCs, trust relationships can be built as well. A cell leader can be elected to represent the members in the cell to communicate with other cells. For security reasons, the cell leader is monitored by its neighbors. When the leader sends and receives aggregated position packets, all the members in the cell will compare the positions in the packets based on their knowledge. By remaining silent, they confirm that the packets have not been altered. Otherwise, they broadcast protest packets against the leader. The other neighbors will put the leader and the protestor vehicle into the question table after receiving the protest packet. Then, the opinion of the other neighbors is taken into account. If the majority of vehicles

regard the leader as malicious, the record of the leader is moved to the distrust table, as discussed [9]. Otherwise, the records sent by the leader are placed in the trust table

6.3.2 Authentication and Confidentiality:

To provide authentication and confidentiality, we propose a geographic location based security mechanism to ensure physical security on top of conventional methods. Messages are encrypted with a geographic location key that specifies a decryption region. This provides physical security because a vehicle has to be physically present in the decryption region to decrypt cipher text encrypted with this geographic location key. As an example, Fig. 6.7 shows a shaded square that is a location-based security region. Sender vehicle a specifies the region, creates the location key, encrypts the message, and sends cipher text to vehicles in this region. Vehicles outside this region such as b, c, d, and e cannot decrypt the message. Only vehicle f can decrypt the message because it is physically inside the decryption region. Since the decryption region can be dynamically specified, attacks are extremely expensive and difficult to mount.

6.4 Enhancing Scalability of Security Schemes

When vehicle population increases in a certain area, not only the scalability of the VC but also the scalability of security schemes becomes a tough problem. In our cloud model, the scalability of the security scheme can be enhanced by a virtual machine division algorithm, a highly scalable algorithm. When the number of access of a virtual machine grows sufficiently large, compared to an empirical threshold, the virtual machines (as a super-VM) will divide itself into multiple sub virtual machines (as sub-VMs). Each virtual machine will obtain the same amount of resources as the original super VM. The middleware of the super VM can randomly forward request to sub virtual machines to load balance. The middleware of the super VM also caches the most recently accessed and frequent information [15]. It caches and executes information such as frequently asked questions (FAQs) and answers. If access from a vehicle hits the FAQ, the middleware directly sends back the answer. If the access misses the FAQ, the middleware then forwards access to a relatively idle VM. This can further reduce the workload of sub-VMs (see Fig. 6.8).

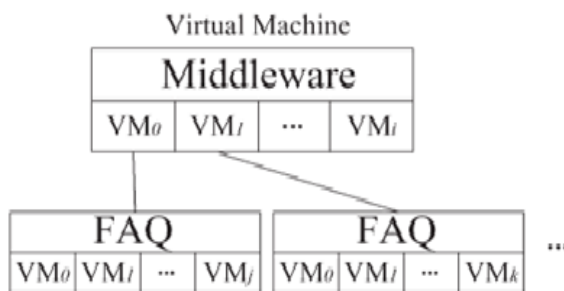


Fig.6.8 Virtual machine can be divided into multilayer's of VMs. Each layer is composed by multiple VMs. The middleware can also be deployed with a cache of frequently ac-

cessed information.

7. ADVANTAGES & DISADVANTAGES

Advantages

- Major applications of VANET include providing safety information, traffic management, toll services, location based services and infotainment.
- One of the major applications of VANET include providing safety related information to avoid collisions, reducing pile up of vehicles after an accident and offering warnings related to state of roads and intersections.
- Affixed with the safety related information are the liability related messages, which would determine which vehicles are present at the site of the accident and later help in fixing responsibility for the accident.
- Collision Avoidance.
- Traffic Optimization.

Disadvantages

- It is cost effective of the system.
- It can not managing time delay.

CONCLUSION

In this seminar, I have addressed the security challenges of a novel perspective of VANETs, i.e., taking VANETs to clouds. I have first introduced the security and privacy challenges that VC computing networks have to face, and I have also addressed possible security solutions. Although some of the solutions can leverage existing security techniques, there are many unique challenges. For example, attackers can physically locate on the same cloud server. The vehicles have high mobility, and the communication is inherently unstable and intermittent. I have provided a directional security scheme to show an appropriate security architecture that handles several, not all, challenges in VCs. We will investigate the brand-new area and design solutions for each individual challenge. Many applications can be developed on VCs. As future work, a specific application will need to analyze and provide security solutions. Extensive work of the security and privacy in VCs will become a complex system and need a systematic and synthetic way to implement intelligent transportation systems. Only with joint efforts and close cooperation among different organizations such as law enforcement, government, the automobile industry, and academics can the VC computing networks provide solid and feasible security and privacy solutions. VANET security is an emerging area. As different VANET protocols and applications are based on different assumptions, a common evaluation framework is needed to compare different security research contributions. Detection of malicious vehicles is still a challenge. Multicast source authentication which essentially guarantees that the received data is sent from the claimed source.

FUTURE WORK

In future work, we will investigate the brand-new area and design solutions for each individual challenge. Many applications can be developed on VCs. As future work, a specific application will need to analyze and provide security solutions. Extensive work of the security and privacy in VCs will become a complex system and need a systematic and synthetic way to implement intelligent transportation systems.

Cost effectiveness of the system

It should be said that implementing our proposed system will lead to many solutions of the security problems that are encountered in VANET. Even the system is costly. So an imperative solution of this system and an effective cost management analysis of this system can be a great future research issue.

Time delay management

VANET is an excellent discovery in terms of safety related information. If the information send later, i.e. after a good amount of time then it will be useless to have such a system. So reducing time delay should be a prime research topic

Using the available technologies such as Wi-Fi, CDMA, GSM

VANET communication uses new protocols. We should think about mixing the communication process with all the existing protocols that are present, such as Wi-fi, CDMA and GSM.

ACKNOWLEDGEMENT

I take this opportunity to express my gratitude and indebtedness to my guide **Prof. R. R. Papalkar**, Computer Science & Engineering department, who is a constant source of guidance and inspiration in preparing this work. I express my sincere gratitude towards our **H.O.D. Prof. A. V. Zade** whose constant help and encouragement helped me to complete my seminar report. I am grateful to **Dr. L. P. Dhamande, Principal** for his encouragement and support.

I am also thankful to all the staff members of Computer Science & Engineering department for helpfully support.

REFERENCES

- [1] Fay Hui: A survey on the characterization of Vehicular Ad Hoc Networks routing solutions .
- [2] D. Huang, S. Misra, G. Xue, and M. Verma, "PACP: An efficient pseudonymous authentication based conditional privacy protocol for vanets," *IEEE Trans. Intell. Transp. Syst.*, vol. 12, no. 3, pp. 736–746, Sep. 2011.
- [3] S. Olariu, I. Khalil, and M. Abuelela, "Taking VANET to the clouds," *Int. J. Pervasive Comput. Commun.*, vol. 7, no. 1, pp. 7–21, 2011.
- [4] L. Li, J. Song, F.-Y. Wang, W. Niehsen, and N. Zheng, "IVS 05: New developments and research trends for intelligent vehicles," *IEEE Intell. Syst.*, vol. 20, no. 4, pp. 10–14, Jul./Aug. 2005.
- [5] S. Arif, S. Olariu, J. Wang, G. Yan, W. Yang, and I. Khalil, "Datacenter at the airport: Reasoning about time-dependent parking lot occupancy," *IEEE Trans. Parallel Distrib. Syst.*, 2012, [Online]. Available: <https://csdl2.computer.org/csdl/trans/td/preprint/ttd2012990021-abs.html>, to be published.
- [6] D. Wen, G. Yan, N. Zheng, L. Shen, and L. Li, "Toward cognitive vehicles," *IEEE Intell. Syst. Mag.*, vol. 26, no. 3, pp. 76–80, May–Jun. 2011.
- [7] S. Arif, S. Olariu, J. Wang, G. Yan, W. Yang, and I. Khalil, "Datacenter at the airport: Reasoning about time-dependent parking lot occupancy," *IEEE Trans. Parallel Distrib. Syst.*, 2012, [Online]. Available: <https://csdl2.computer.org/csdl/trans/td/preprint/ttd2012990021-abs.html>, to be published.
- [8] R. Hasan, *Cloud Security*. [Online]. Available: <http://www.ragibhasan.com/research/cloudsec.html>
- [9] S. Olariu, M. Eltoweissy, and M. Younis, "Toward autonomous vehicular clouds," *ICST Trans. Mobile Commun. Comput.*, vol. 11, no. 7–9, pp. 1–11, Jul.–Sep. 2011.
- [10] SIRIT-Technologies, White paper. DSRC technology and the DSRC industry consortium (DIC) prototype team.
- [11] L. Li, J. Song, F.-Y. Wang, W. Niehsen, and N. Zheng, "IVS 05: New developments and research trends for intelligent vehicles," *IEEE Intell. Syst.*, vol. 20, no. 4, pp. 10–14, Jul./Aug. 2005.
- [12] G. Yan and S. Olariu, "A probabilistic analysis of link duration in vehicular ad hoc networks," *IEEE Trans. Intell. Transp. Syst.*, vol. 12, no. 4, pp. 1227–1236, Dec. 2011.
- [13] H. Xie, L. Kulik, and E. Tanin, "Privacy-aware traffic monitoring," *IEEE Trans.*
- [14] T. Ristenpart, E. Tromer, H. Shacham, and S. Savage, "Hey, you, get off of my cloud: Exploring information leakage in third-party compute clouds," in *Proc. 16th ACM Conf. CCS*, 2009, pp. 199–212.
- [15] A. Friedman and D. West, "Privacy and security in cloud computing," *Center for Technology Innovation: Issues in Technology Innovation*, no. 3, pp. 1–11, Oct. 2010.
- [16] J. A. Blackley, J. Peltier, and T. R. Peltier, *Information Security Fundamentals*. New York: Auerbach, 2004.
- [17] N. Santos, K. P. Gummadi, and R. Rodrigues, "Toward trusted cloud computing," in *Proc. HotCloud*, Jun. 2009
- [18] T. Garfinkel, B. Pfaff, J. Chow, M. Rosenblum, and D. B. Terry, "Virtual machine-based platform for trusted computing," in *Proc. ACM SOSP*, 2003, pp. 193–206.
- [19] S. Berger, R. Cáceres, K. A. Goldman, R. Perez, R. Sailer, and L. van Doorn, "VTPM: Virtualizing the trusted platform module," in *Proc. 15th Conf. USENIX Sec. Symp.*, Berkeley, CA, 2006, pp. 305–320.
- [20] D. G. Murray, G. Milos, and S. Hand, "Improving XEN security through disaggregation," in *Proc. 4th ACM SIGPLAN/SIGOPS Int. Conf. VEE*, New York, 2008.
- [21] F. J. Krauthem, "Private virtual infrastructure for cloud computing," in *Proc. Conf. Hot Topics Cloud Comput.*, 2009, pp. 1–5
- [22] Q. Wang, C. Wang, J. Li, K. Ren, and W. Lou, "Enabling public verifiability and data dynamics for storage security in cloud computing," in *Proc. 14th ESORICS*, 2009.
- [23] M. Jensen, J. Schwenk, N. Gruschka, and L. L. Iacono, "On technical security issues in cloud computing," in *Proc. IEEE Int. Conf. Cloud Comput.*, 2009, pp. 109–116